



an ACK request message to notify that the 200 OK response is corrected received. Since the ACK message is used only for confirmation, we define the call setup latency as the interval from the time instant the INVITE message is sent to the time instant the 200 OK message is received.

Since the delay in wired links,  $t_{wired}$ , is fairly stable, it is assumed that the call setup procedure is delayed only due to wireless link delay. For WLAN,  $N$  saturated MNs (i.e., each MN always has a packet to send) are assumed to share the WLAN channel and the transmission failures are only due to collisions. We consider the IEEE 802.11 distributed coordination function (DCF) [1] with a basic access mode for media access control because request-to-send (RTS)/clear-to-send (CTS) is not quite effective in infrastructured WLANs and it is disabled in most products available in the current market.

We assume that user datagram protocol (UDP) is used as a transport protocol for SIP messages. Since UDP does not support reliable transmissions, the UAC performs end-to-end (E2E) retransmissions based on an exponential backoff algorithm [2]. The initial backoff timer  $T_{Init}$  is typically set to 500 msec. After the UAC sends an INVITE message, the UAC retransmits the INVITE message at most for 32 seconds or until it receives a response. As a result, the UAC can retransmit an INVITE message at most 6 times, that is, the number of total transmissions is 7. Similarly, the UAS retransmits a 200 OK message at most for 32 seconds or until it receives an ACK message. However, unlike the UAC, the UAS keeps the maximum backoff timer at 4 seconds. Therefore, retransmissions of a 200 OK message occur at most 10 times, i.e., at 0.5, 1.5, 3.5, 7.5, 11.5, 15.5, 19.5, 23.5, 27.5, and 31.5 seconds after the transmission of the first 200 OK message. Independently from E2E retransmissions, MAC layer retransmissions are performed in IEEE 802.11 WLANs. That is, when a packet<sup>2</sup> transmission fails due to collision, the packet is retransmitted until it is successfully delivered or up to  $m$  times.

### III. VOIP CALL SETUP LATENCY ANALYSIS

Let  $\varepsilon$  be the probability that a SIP message is lost over a WLAN link after  $m$  MAC layer retransmissions. Since a packet transmission at each backoff stage is independent,  $\varepsilon = p^{m+1}$  where  $p$  is the collision probability for a packet transmission.  $p$  can be obtained from [9] by an iterative method. Let  $R$  be the total number of E2E transmissions. Then,  $R$  is 7 and 11 for the INVITE and 200 OK messages, respectively. By [1], the number of backoff slots in the  $j$ th stage is uniformly selected in  $[0, W_j - 1]$  where  $W_j = 2^j W_0$  and  $W_0$  is the minimum contention window size. Therefore, the probability that the number of chosen backoff slots is  $k$  at the  $j$ th stage is  $1/(W_j - 1)$ .

Let  $\theta_X(i, j, k)$  be the probability that a transmission of a SIP message  $X$  is successful at the  $i$ th E2E transmission and the  $j$ th backoff, and the chosen backoff slot length is  $k$  ( $1 \leq$

$i \leq R$ ,  $0 \leq j \leq m$ ,  $1 \leq k \leq W_j - 1$ ). Then,  $\theta_X(i, j, k)$  can be obtained from

$$\theta_X(i, j, k) = \frac{1}{1 - \varepsilon^R} \frac{\varepsilon^{i-1} p^j (1 - p)}{W_j - 1}, \quad (1)$$

where  $1/(1 - \varepsilon^R)$  is the normalized factor that is required because we consider only successful transmissions.

Let  $l_X(i, j, k)$  be the transmission time when the counters for E2E transmission and MAC layer transmission for a transmission of a SIP message  $X$  are  $i$  and  $j$ , respectively, and the randomly chosen contention window size is  $k$  ( $1 \leq i \leq R$ ,  $0 \leq j \leq m$ ,  $1 \leq k \leq W_j - 1$ ). Then,  $l_X(i, j, k)$  is given by (2), where  $Tr(n)$  is the value of the E2E retransmission timer when a packet is successfully transmitted at the  $n$ th transmission.  $Tr(n)$  for the INVITE and 200 OK messages are given by  $2^{n-2}T_{Init}$  and  $\min\{2^{n-2}T_{Init}, 4\}$ , respectively, where  $T_{Init}$  is the initial E2E timeout value. In (2),  $E[slot]$  is the average slot length defined as the time interval between two consecutive backoff counter decrements in [9], and it can be computed as

$$E[slot] = (1 - P_{tr})\sigma + P_{tr}P_S T_S + P_{tr}(1 - P_S)T_C,$$

where  $P_{tr} = 1 - (1 - \tau)^n$  and  $P_S = \frac{n\tau(1-\tau)^{n-1}}{1-(1-\tau)^n}$ .  $T_S$  and  $T_C$  are the time durations that the channel is sensed busy during a successful frame transmission and a collision, respectively.  $T_S$  and  $T_C$  are given by

$$T_S = DIFS + H + P + \delta + SIFS + ACK + \delta \quad \text{and}$$

$$T_C = DIFS + H + P + SIFS + ACK,$$

where  $\delta$  is the propagation delay.  $DIFS$  and  $SIFS$  represent DCF inter frame space and small inter frame space, respectively.  $H$ ,  $P$ ,  $ACK$  are the transmission time for the header, payload, and ACK frame, respectively.

Then, the average transmission time of a SIP message  $X$  is

$$L(X) = \sum_{i=1}^R \sum_{j=0}^m \sum_{k=1}^{W_j-1} \theta_X(i, j, k) \cdot l_X(i, j, k). \quad (3)$$

The average call setup latency can be represented as

$$S = L(INVITE) + L(200OK), \quad (4)$$

where  $L(INVITE)$  and  $L(200OK)$  are delivery latency for the INVITE and 200 OK messages, respectively.

### IV. SIMULATION RESULTS

To validate the analytical model, we have carried out simulations using ns-2 simulator [10]. In simulations, the wired link delay is fixed at 20 msec. The retransmission limit is set to 3 and the data rate is 1 Mbps. Note that  $m = 3$  is less than the value in the IEEE 802.11 specification. This is because we focus on time-sensitive VoIP applications and thus a larger  $m$  is not appropriate due to a long end-to-end delay and delay jitter even though it can reduce the packet loss rate. The lengths (including UDP and IP headers) of INVITE and 200 OK messages are 788 and 488 bytes, respectively.

Figure 2 shows the call setup latency as a function of  $N$  under different  $T_{Init}$ . It can be seen that the call setup latency increases with the increase in  $N$ . Also, the latency

<sup>2</sup>Since the maximum protocol data unit in IEEE 802.11 MAC layer is sufficiently large, we do not consider link layer fragmentation and therefore the term *packet* is used for both a protocol data unit (PDU) both in the data link layer and in the transport layer.

$$l_X(i, j, k) = \begin{cases} \sum_{n=2}^i Tr(n) + t_{wired} + \sum_{m=0}^{j-1} \frac{W_m-1}{2} E[slot] + kE[slot], & i > 1, j \geq 1 \\ \sum_{n=2}^i Tr(n) + t_{wired} + \frac{W_0-1}{2} E[slot] + kE[slot], & i > 1, j = 0 \\ t_{wired} + \sum_{m=0}^{j-1} \frac{W_m-1}{2} E[slot] + kE[slot], & i = 1, j \geq 1 \\ t_{wired} + \frac{W_0-1}{2} E[slot] + kE[slot], & i = 1, j = 0 \end{cases} \quad (2)$$

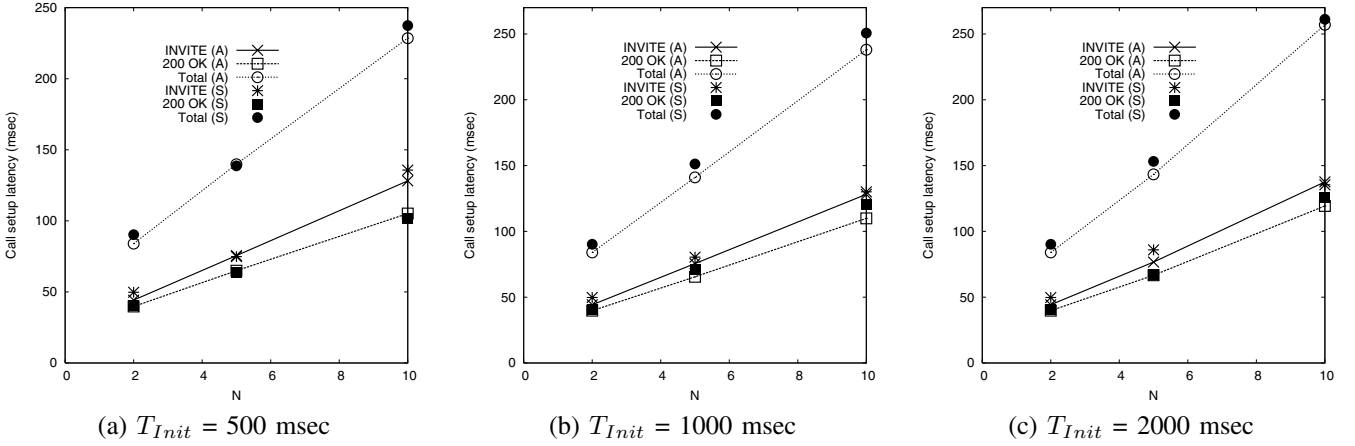


Fig. 2. Call setup latency vs.  $N$  (A: Analysis, S: Simulation).

for the INVITE message is larger than that for the 200 OK message because the INVITE message is longer than the 200 OK message. From Figures 2, it can be found that some discrepancy between simulation and analytical results especially when there are many contending mobile nodes (i.e.,  $N$  is large). This discrepancy can be explained by two reasons: 1) the AP queueing delay is not considered in our model. However, the AP queueing delay cannot be neglected when  $N$  is large; 2) unnecessary retransmissions can occur for a large  $N$  if a packet is retransmitted before the original packet is completely served at the AP queue due to long channel contention time. Since the unnecessary retransmission is not taken account in our model, larger discrepancy can be observed in the situation. Actually, the AP queueing delay can be mitigated by increasing the initial E2E retransmission timer,  $T_{Init}$ . Consequently, as shown in Figure 2, the discrepancy between simulation and analytical results decreases as  $T_{Init}$  increases. It can be also seen that the average call setup latency also increases with the increase of  $T_{Init}$ . Therefore, an adaptive setting for  $T_{Init}$  depending on network conditions needs to be devised to optimize the call setup latency.

## V. CONCLUSION

In this letter, we have analyzed the call setup latency in a SIP-based VoWLAN system. Extensive simulations have been carried out to validate the analytical model. Simulation and analytical results demonstrate that the call setup latency is sensitive to the number of mobile nodes in a WLAN, and  $T_{Init}$  should be carefully chosen to avoid unnecessary end-to-

end retransmissions while minimizing the call setup latency. In our future work, we will exploit the call setup blocking probability, which is defined as the probability the call setup latency exceeds an acceptable bound, and extend the analytical model for wireless mesh networks with multi-hop wireless links.

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