Zero-Knowledge Middleboxes

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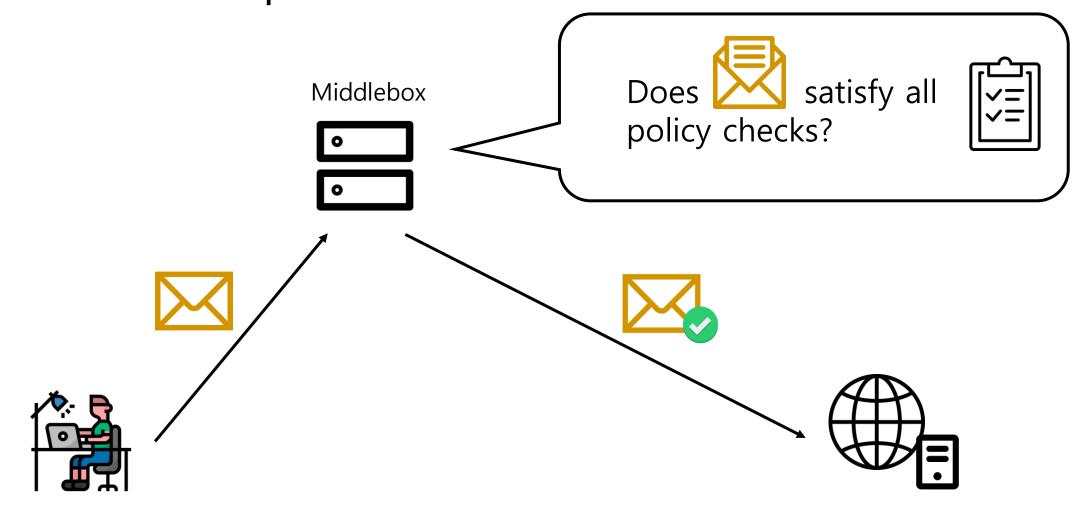
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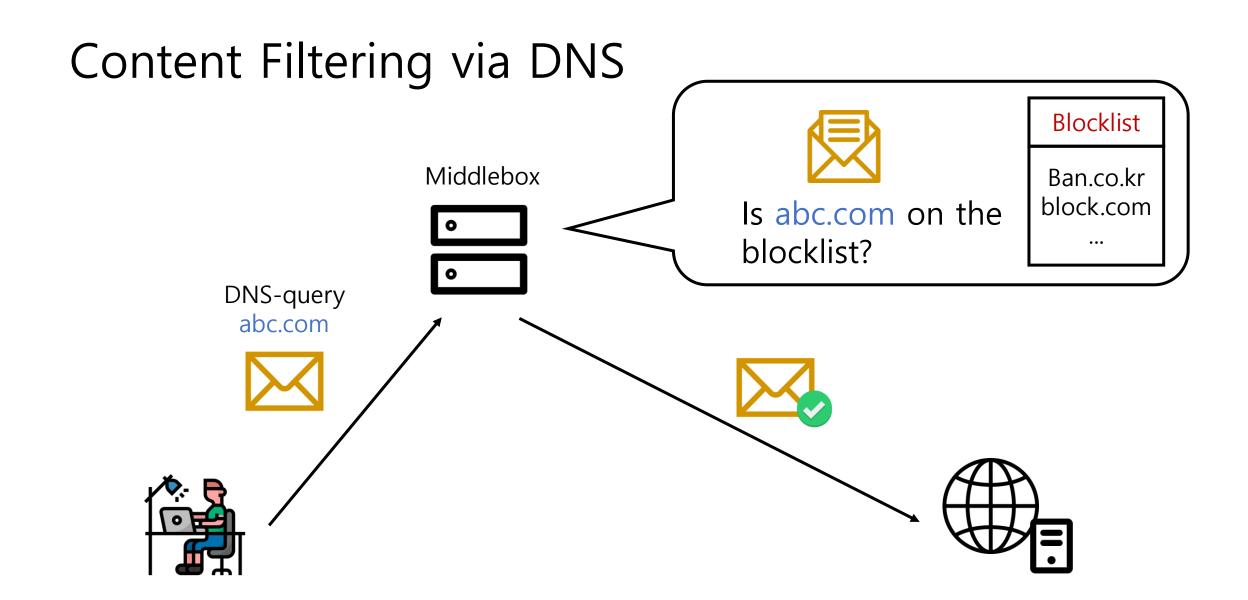
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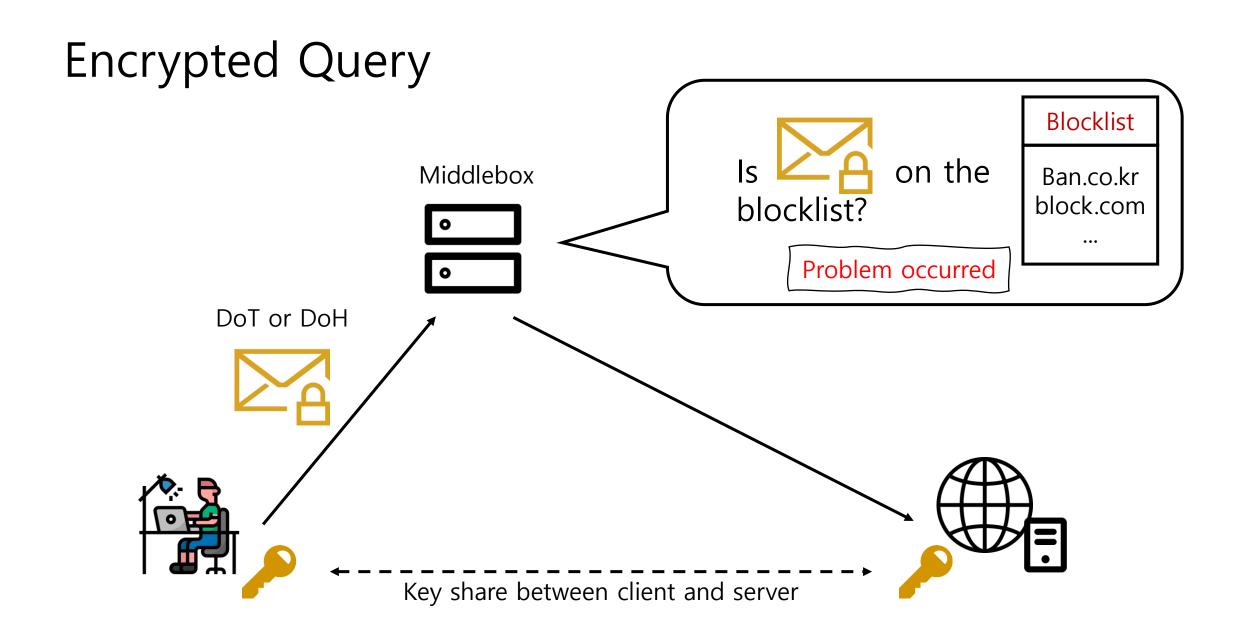
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Middlebox inspects traffic







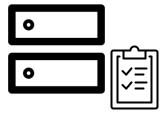
Motivation

Client



Want to keep own **Privacy**

Middlebox



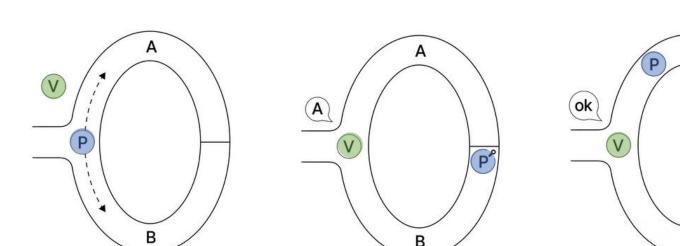
Want to imforce own **Policy**

Compatible Solution

Zero-Knowledge Proofs

Background – Zero Knowledge Proofs

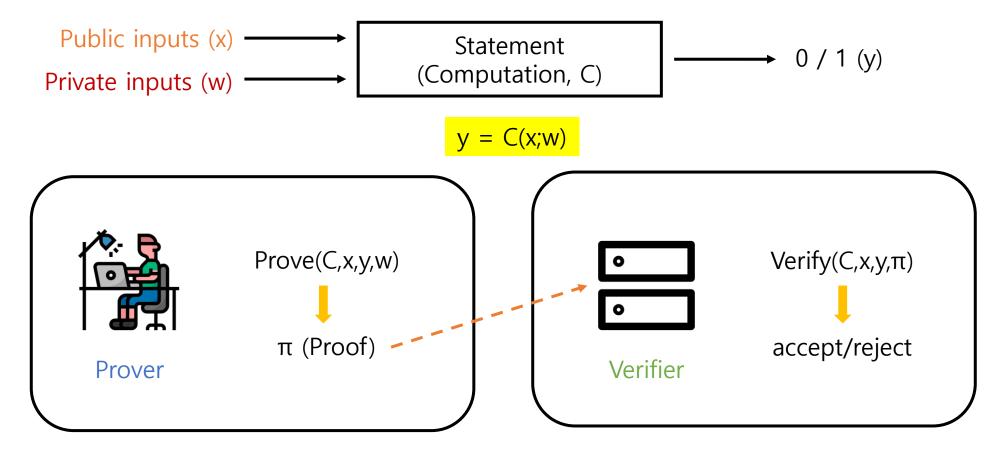
- Procedure in cryptography that ensures that when someone proves to another that a statement is true, they do not reveal anything except whether the statement is true or false
- The way to prove the <u>validity</u> of information without revealing any information
- Prover prove to the Verifier that it knows the Secret without revealing it



Probability

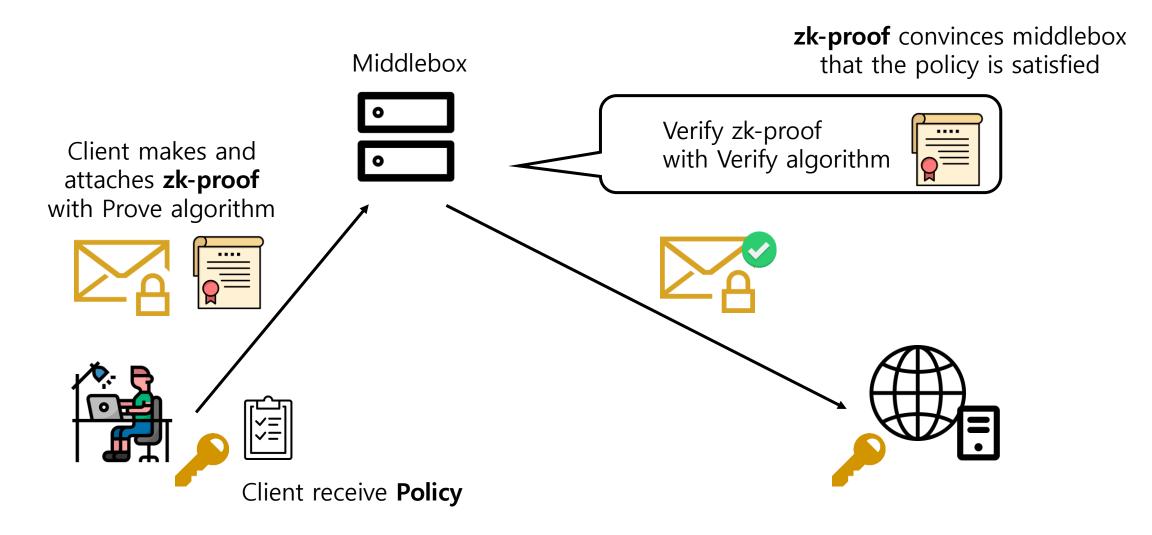
1 time: 50% 20 times: <0.0001%

Background – Zero Knowledge Proofs (cont.)



QAP-based proof protocol Groth16 (zkSNARK) which is a non-interactive zero knowledge proof (NIZK)

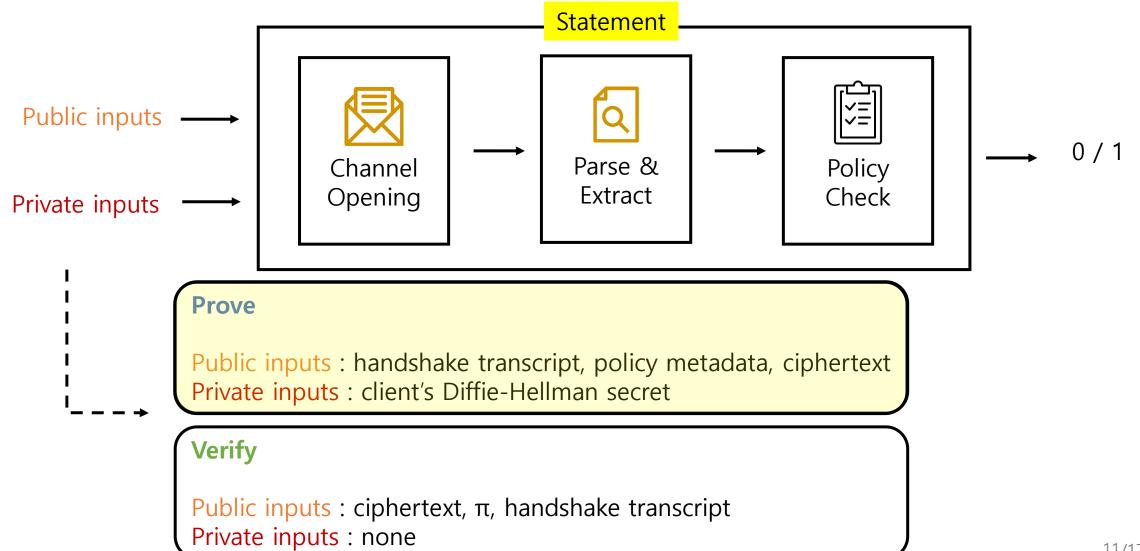
The ZKMB Framework



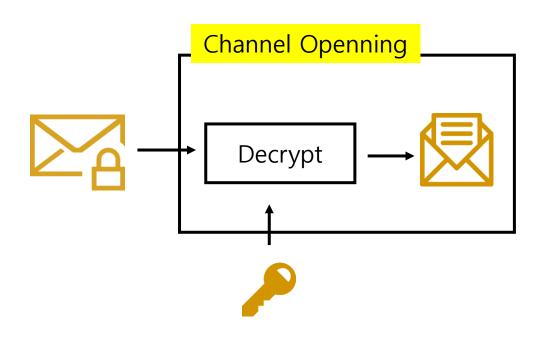
The challenge with network messages

- Many practical ZK protocols (e.g., SNARKs) are made
 - They require representing the statement as an arithmetic circuit over a field
- Legacy symmetric-key functions like AES, ChaCha, SHA are very inefficient as circuits
 - Thus, they **must design proof statements** for this case
- Of course, No weakened encryption or privacy guarantees of TLS 1.3 and no server-side modification

ZKMB proof statements

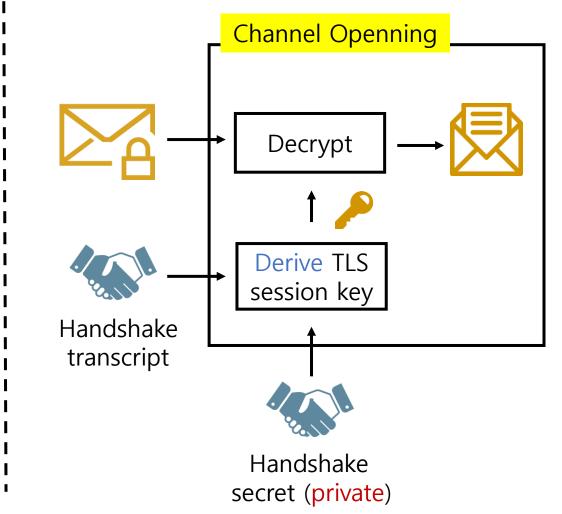


Channel Opening



Challenge

Prover may use a different key because TLS 1.3 doesn't support any "key-committing" authenticated encryption ciphers

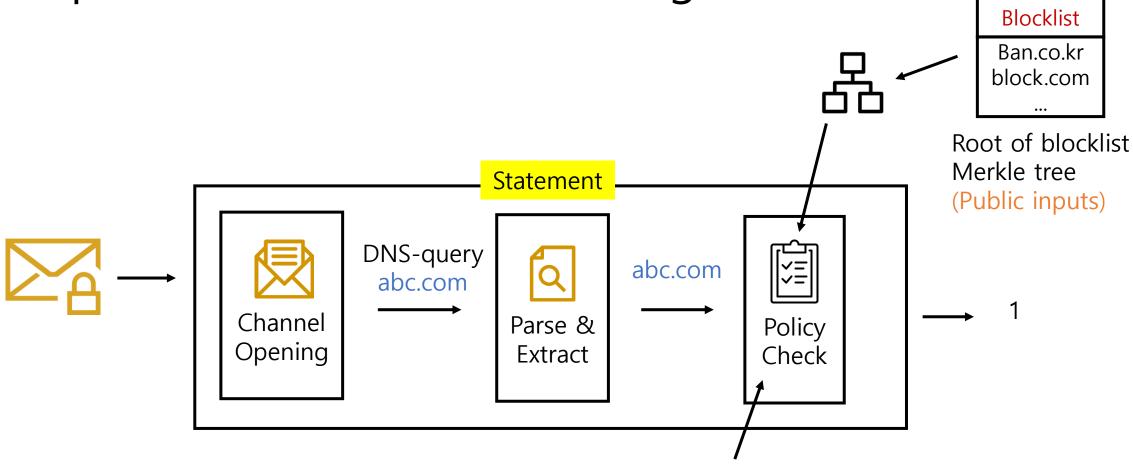


Parse & Extract and Policy Check

- Translation between the network protocol wire format of message and the input format of the policy check
 - Extract the policy-relevant substring of a network packet, while checking that some syntactic requirements
 - Operate differently depending on whether it's DoH or DoT, because they have different format of message

• Check that extracted message satisfies the policy in the policy metadata

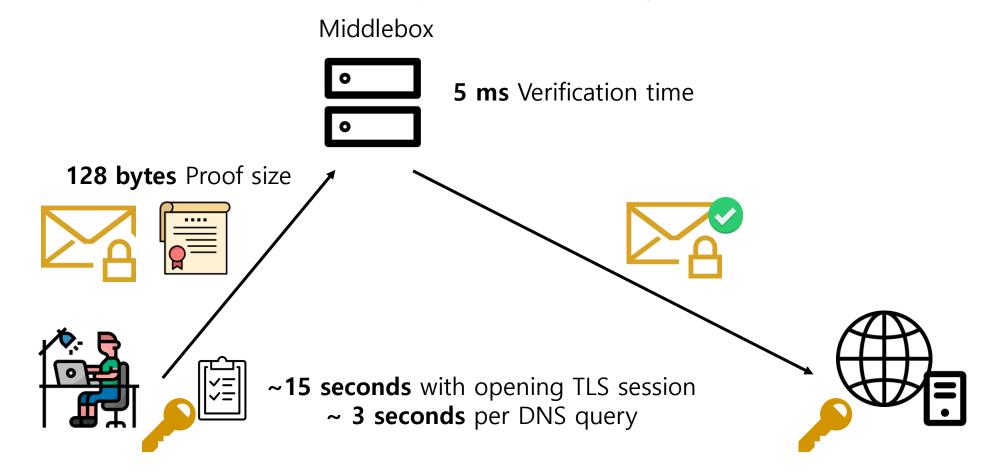
Implementation – DNS filtering



Merkle non-membership proof of URL (Private inputs)

Experiment Result

• Implemented with xJsnark library and Groth16 algorithm



Conclusion

- ZKMB is **compatible solution** for clients and middleboxes
 - Clients prove policy-compliance using zk-proof with keeping privacy

- Modular policy check
 - DNS filtering, HTTP firewall, ...

- Future work
 - Low delay
 - Protocol-specific optimization
 - Not only in local network, but also in cloud network

Thank you for listening